

Challenges in High Performance Computing

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Challenges in High Performance Computing

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Overview

- Demands for HPC Storage: Parallel Filesystems
 - Lustre
 - Panasas
 - Parallel NFS
- Demands of SMP applications: much RAM, many cores
 - vSMP
- Demands for easy cluster deployment and management
 - Microsoft HPC Server 2008
 - Zaratustra

1. Demands for HPC Storage

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Storage Demands in HPC

- need for **computing power**
 - due to need to run larger and more accurate models
 - more CPUs, more cores, more nodes, more RAM
- need for **network performance**
 - more highly parallelized jobs
 - high-speed interconnects (10GbE, InfiniBand,...)

→ **massive explosion of data sets**

→ demand for

- **large storage capacity**
- **high bandwidth**
- **low latency**

Deficiencies of Today's Solutions

- most widespread solution: **single NFS server**
 - does not scale: NFS head is bottleneck
 - „high-speed“ NFS server will be bottleneck by tomorrow
- „**clustered NFS**“: problematic
 - either head-to-head synchronization limits scalability
 - or manual partitioning of global namespace is cumbersome
 - NFS is not suitable for dynamical load balancing (inherent state)

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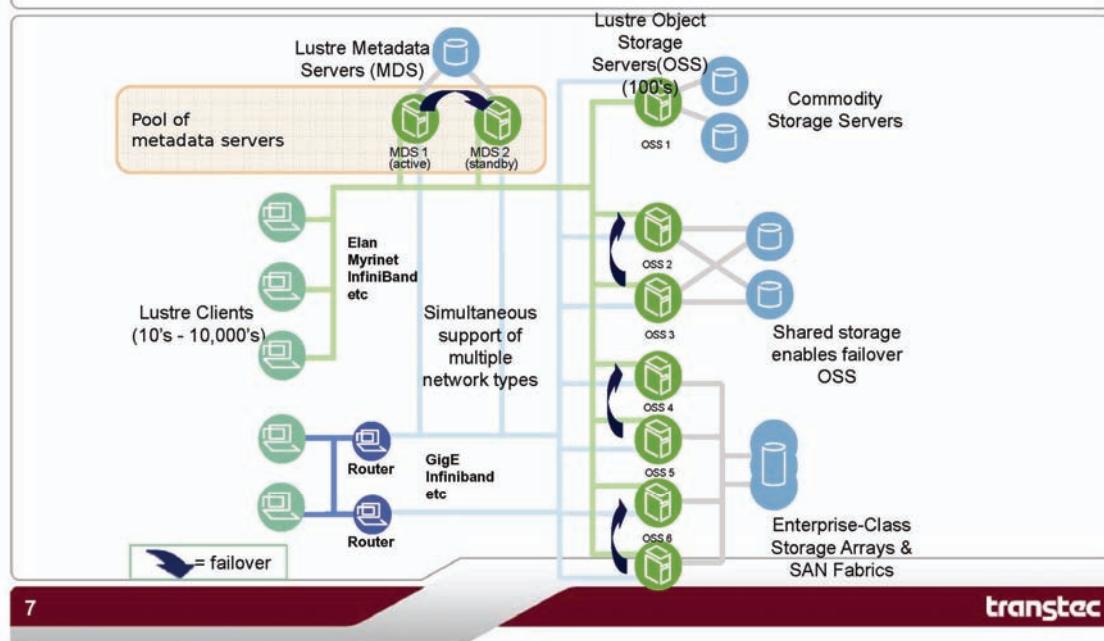
Parallel File Systems

- major features:
 - **global namespace** eases filesystem management and job flow
 - **scalable** capacities and bandwidths
- **distributed vs. cluster vs. parallel filesystem:**
 - no shared storage → many-to-many access to data
- **proprietary solutions** already there:
 - IBM's GPFS
 - SGI's CXFS
 - Panasas' ActiveScale Filesystem (PanFS)
 - EMC's Celerra MPFS/MPFSi (pka High Road)
 - Lustre, PVFS2, ...

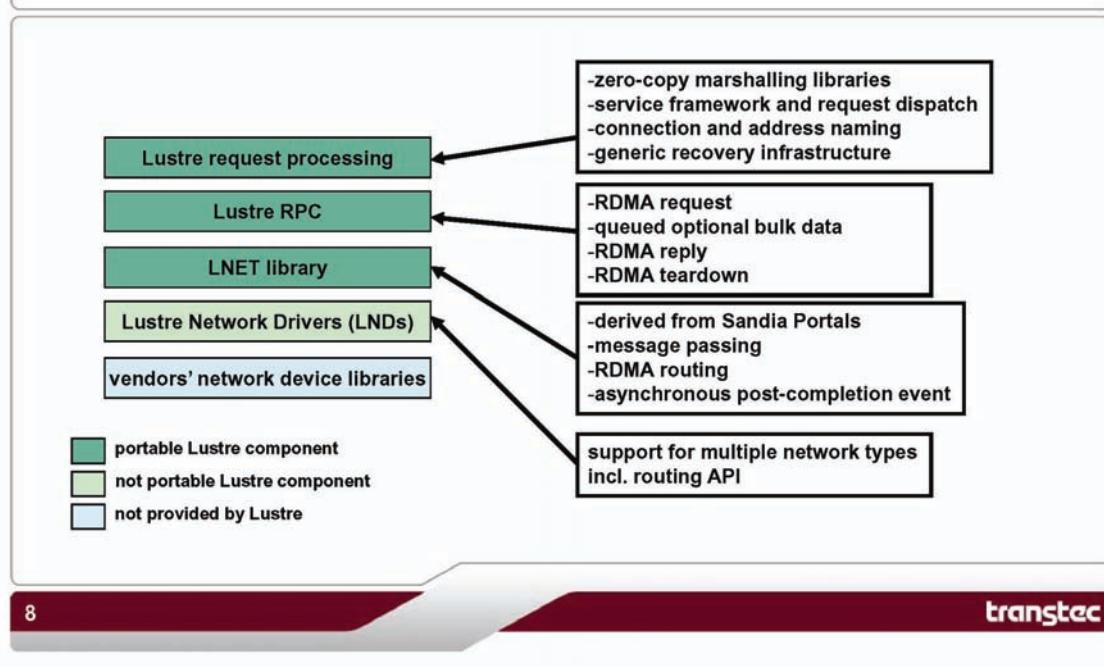
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Lustre Overview



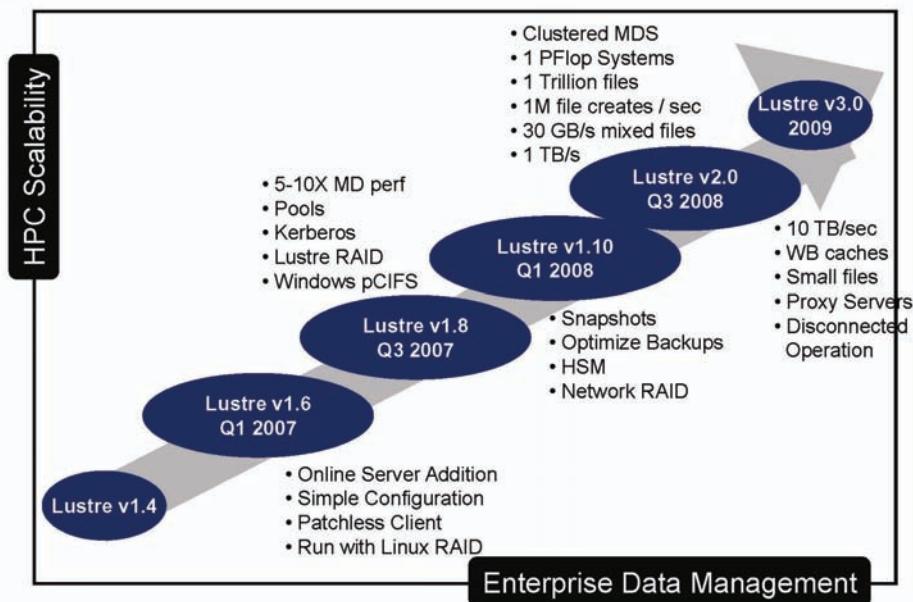
Lustre – Modular Network Implementation



Lustre Limits

- maximum stripe count: 160
maximum number of OSTs: 1020
maximum number of clients: 32768 (tested: 22000)
maximum size of filesystem: 8 TB * #OSTs

Intergalactic Strategy



NFS as a Standard

- need for **OS independent, interoperable, standardized** solution
→ **NFS is the ONLY standard!**
- **standards are good, because...**
 - they **protect end user investment** in technology
 - they **ensure a base level of interoperability**
 - while at the same time **provide choice** among products
 - commonality leads to **less training, simpler deployment, higher acceptance...**

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NFS 4.1 and Parallel NFS (pNFS)

- **NFS 4.1:** idea to use SAN FS architecture for NFS originally from Gary Grider (LANL) and Lee Ward (Sandia)
- development driven by Panasas, Netapp, Sun, EMC, IBM, UMich/CITI
- folded into NFSv4 minor version NFSv4.1 in 2006
- future **internet standard**
(current draft 21: <http://www.ietf.org/internet-drafts/draft-ietf-nfsv4-minorversion1-21.txt>)
- major changes to NFS 4:
 - sessions
 - directory delegations
 - pNFS (optional feature)
- standardization expected some time in 2009

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Parallel NFS (pNFS): Generic Architecture

- **separation of metadata path and data path (out-of-band global namespace)**
- built for **interoperability** and **backwards-compatibility**
- **flexible design allows for different storage implementations (layouts)**



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What pNFS Does NOT Give You

- **improved cache consistency**
 - NFS has **open-to-close consistency**
- **perfect POSIX semantics** in a distributed file system
- **clustered metadata**
 - though a mechanism for this is not precluded

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Parallel NFS (pNFS): Different Layout Formats

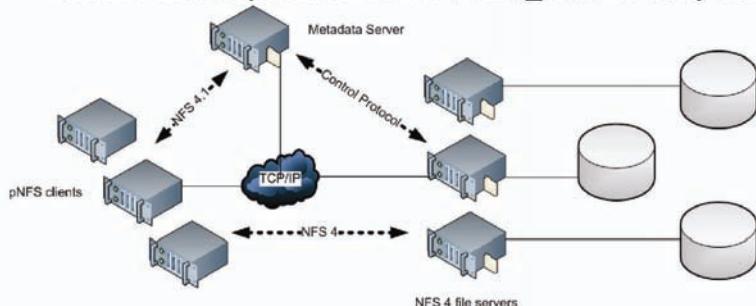
- a **layout** describes the location of file data, containing a list of device IDs and striping information
- possession of a layout grants access to storage devices, resp. files
- **file-based layout** (part of NFS 4.1/pNFS standard)
- **block-based layout**: <http://www.ietf.org/internet-drafts/draft-ietf-nfsv4-pnfs-block-08.txt>
- **object-based layout**: <http://www.ietf.org/internet-drafts/draft-ietf-nfsv4-pnfs-obj-07.txt>
- PVFS2 layout
- GPFS layout
- ...

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pNFS: File Layout

- only storage access protocol directly specified in NFS 4.1 standard
- significantly co-designed by NetApp, Sun, IBM and others
- file layout simple, may be **heavily cached** by clients
- access control possible via RPCSEC_GSS security flavor

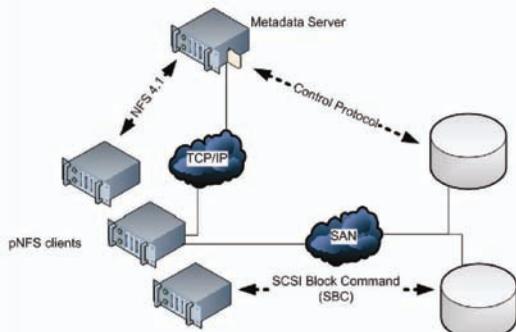


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pNFS: Block Layout

- highly influenced by EMC's design of **Multi-Path File System MPFS(i) (pka High Road)**
- block layout uses **volume identifiers, block offsets and extents**
- secure authorization with **host granularity only**, file-level security cannot be enforced by storage devices
→ **clients must be trusted** (fundamental NFS problem ever since)

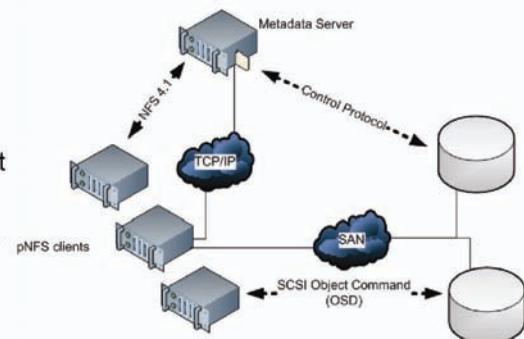


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pNFS: Object Layout

- Panasas' contribution, based on **NASD** design (Network-Attached Secure Disk) developed at Carnegie Mellon University, later evolved into forthcoming **SCSI OSD** standard (**object-based storage device**)
- layout uses **SCSI object command set**
- space management** built into devices
- designed for **secure access and high-performance data replication**
- cryptographically secured credentials ("capabilities") needed to access storage devices



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pNFS: The Current State

- **Linux:** file layout, based on PVFS2 / based on NFS 4
- **OpenSolaris:** file (NFS 4) / object (OSD-1) layout driver will be completed soon, patches available: <http://opensolaris.org/os/project/nfsv41/> <http://opensolaris.org/os/project/osd/>
- **Netapp:** file layout, based on NFS 4
- **IBM:** file layout, based on GPFS
- **EMC:** block layout, based on MPFS(i)
- **Panasas:** object layout, based on ActiveScale PanFS
- **Carnegie Mellon University:** performance and correctness testing

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Weblinks Lustre

- Lustre Wiki: <http://wiki.lustre.org>
- Lustre Manual: <http://manual.lustre.org>
- Lustre Knowledge Base:
<https://bugzilla.lustre.org/showdependencytree.cgi?id=2374>

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Weblinks pNFS

- **NASD: Network Attached Secure Disks:** <http://www.pdl.cmu.edu/NASD/>
- **Panasas:** www.panasas.com
- **EMC Celerra Multi-Path File System:**
<http://www.emc.com/products/detail/software/celerra-multipath-file-system.htm>
- **pNFS Information Portal:** <http://www.pnfs.com>
- **NFSv4 Status Pages:** <http://tools.ietf.org/wg/nfsv4>
- **Object-Based Storage Devices (now INCITS 400-2004):**
<http://www.t10.org/ftp/t10/drafts/osd/osd-r10.pdf>
Object-Based Storage Devices V2:
<http://www.t10.org/ftp/t10/drafts/osd2/osd2r03.pdf>
- **Eisler's NFS Blog:** http://blogs.netapp.com/eislernfs_blog
- **NFSv4.1 Bakeathon at OpenSolaris.org:**
http://opensolaris.org/os/project/nfsv41/nfsv41_bakeathon/

2. Demands of SMP Applications

Demands of SMP Applications

- **large memory** for applications
 - to enable workloads that cannot be run otherwise
- **shared memory** coupled with **many CPU cores**
 - to allow multi-threaded applications to benefit from shared-memory systems
- **BUT:** proprietary SMP machines (pka “mainframes”) very expensive and not compatible!

→ **vSMP™ (“versatile SMP”) technology:**

- aggregates multiple off-the-shelf x86 server boards into one virtual x86 system

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vSMP Foundation™ Technology

Multiple off-the-shelf x86 boards, with processors and memory
Processors speed/amount or memory amount doesn't have to be same across all boards



InfiniBand HCAs, cables and switch



vSMP Foundation™ Devices

The flash-devices plug into the boards and used as bootable device. vSMP Foundation is booted to present an aggregate coherent view to the OS



High-end x86 system, based on standard x86 components

Standalone

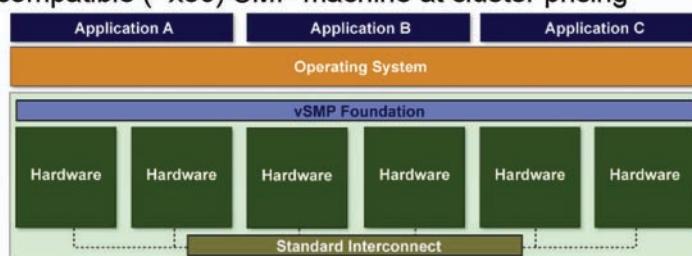
Embedded

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vSMP Foundation™ Technology

- one system to manage: fewer, larger nodes means less cluster management overhead
 - single operating system instance
 - avoid cluster file systems
 - complexities of InfiniBand hidden
- compatible (=x86) SMP machine at cluster pricing



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vSMP Foundation™ Technology: Behind the Scenes

One System

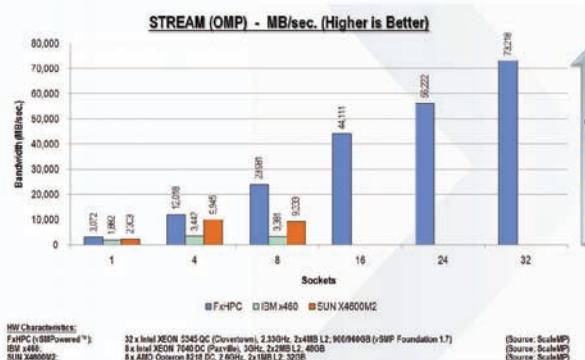
- Software interception engine creates a uniform execution environment
- vSMP Foundation creates the relevant BIOS environment to present the OS (and the SW stack above it) as single coherent system

Coherent Memory

- vSMP Foundation maintains cache coherency between boards
- Multiple concurrent memory coherency mechanisms, on a per-block basis, based on real-time memory activity access pattern
- Leverage board local-memory for caching

Shared I/O

- vSMP exposes all available I/O resources to the OS in a unified PCI hierarchy
- No need for cluster file systems



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vSMP Compatible Systems (1)

- **system capabilities:**
 - max. boards: 16
 - max. memory per board: 4 GB / 64 GB
 - max. processors per board: 1 / 4
 - max. cores per board: 1 / 8
 - max. memory: 1 TB
 - max. processors: 64
 - max. cores: 128
- **fault-tolerant:** supported, minimum boards: 3

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vSMP Compatible Systems (2)

- **supported hardware:**
 - processors: Intel Xeon
 - systems (excerpt):
 - discrete server systems:
 - Supermicro: 6015TW-INF/B / 6015TW-INF/V
 - blade-server systems:
 - Supermicro: SuperBlade: SBI-7125W-S6
- **supported operating systems:**
 - Linux OS, kernel level 2.6.11 or later
 - Red Hat Enterprise Linux 4/5 (RHEL4-5)
 - SUSE Linux Enterprise Server 10 (SLES10)
 - Fedora Core 4/5/6/7 (FC4-7)
 - OpenSUSE 10 or later



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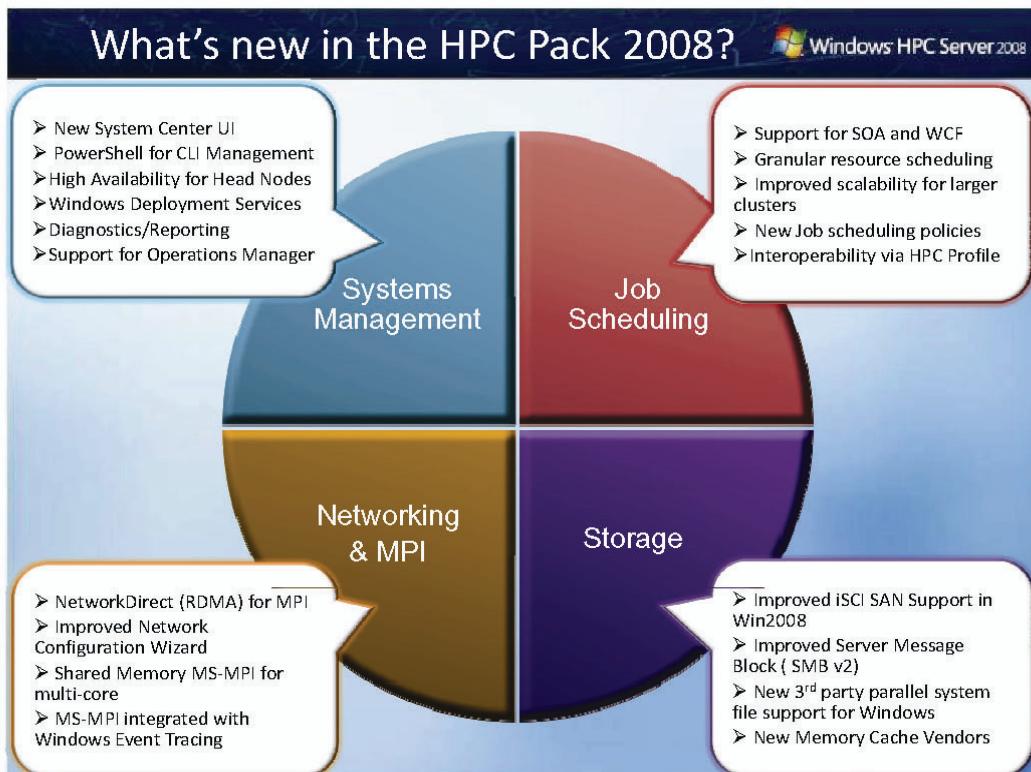
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3. Demands for Easy Cluster Deployment and Management

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Demands for Easy Cluster Deployment

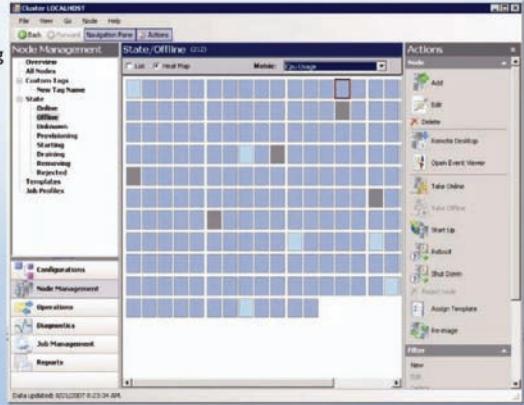
- **WANTED: versatility:**
 - clusters may be partitioned into different entities (different CPUs, different amount of RAM or disk space, various applications etc.)
 - different purchase generations and extensions during lifecycle lead to inhomogeneities in hardware
 - stateless vs. diskless vs. fully installed nodes
 - different OSses
- **BUT also simplicity:**
 - single point of administration
 - abstraction of node varieties from management software
- **AND: integrability:**
 - cluster should not be an „island“ where everything is different



Systems Management - Overview

Windows HPC Server 2008

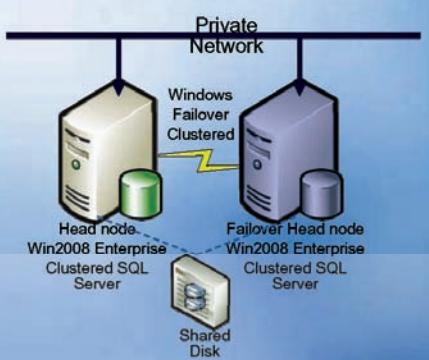
- Simple to setup and manage in a familiar environment
 - Turnkey cluster solutions through OEMs
 - Simplify system and application deployment
 - Base images, patches, drivers, applications
- Focus on ease of management
 - Comprehensive diagnostics, troubleshooting and monitoring
 - Familiar, flexible and “pivotal” management interface
 - Equivalent command line support for unattended management
- Scale up
 - Scale deployment, administration, infrastructure
 - Head node failover
 - Cluster usage reporting
 - Compute node filtering
- Better integration with enterprise management
 - Patch Management
 - System Center Operations Management
 - PowerShell
 - Windows 2008 high Availability Services



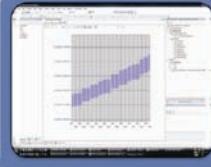
Head Node High Availability

Windows HPC Server 2008

- Eliminates single point of failure with support for high availability
- Requires Windows Server 2008 Enterprise *Failover Clustering Services*
 - Next generation of cluster services
 - Major improvement in configuration validation and management
- HPC Pack Includes
 - Setup integration with Failover Clustering Services
 - Head Node and Failover Node set up with SQL Failover Cluster
 - Job Scheduler services failover
 - Management console linked to Windows Server Failover Management console

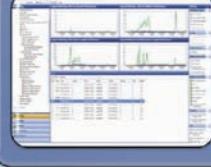


System Center Operations Manager for HPC Windows HPC Server 2008



A more productive HPC environment

- Canned reports for end-user perspective monitoring
- Security logs analysis and reporting



Scalable Monitoring

- Monitor apps running in a scale out, distributed environment
- Scale using tiered management servers
- Agent-less Monitoring



Increased Efficiency and Control

- More secure by design
- Integration with Active Directory
- Extended solution with Management Packs

HPC++

Windows HPC Server 2008

Job Scheduling

HPC++

What's new in Job Scheduling

Windows HPC Server 2008

Broader application support

Expanded Job Policies

Support for Job Templates

Improve interoperability with mixed IT infrastructure

HPC++

Broader Application Support

Windows HPC Server 2008

2003 (focusing on batch jobs)			2008 (focusing on Interactive jobs)		
Engineering Applications Structural Analysis Crash Simulation	Oil & Gas Applications Reservoir simulation Seismic Processing	Life Science Applications Structural Analysis Crash Simulation	Financial Services Portfolio analysis Risk analysis Compliance Actual	Excel Pricing Modeling	Interactive Cluster Applications Your applications here

Job Scheduler

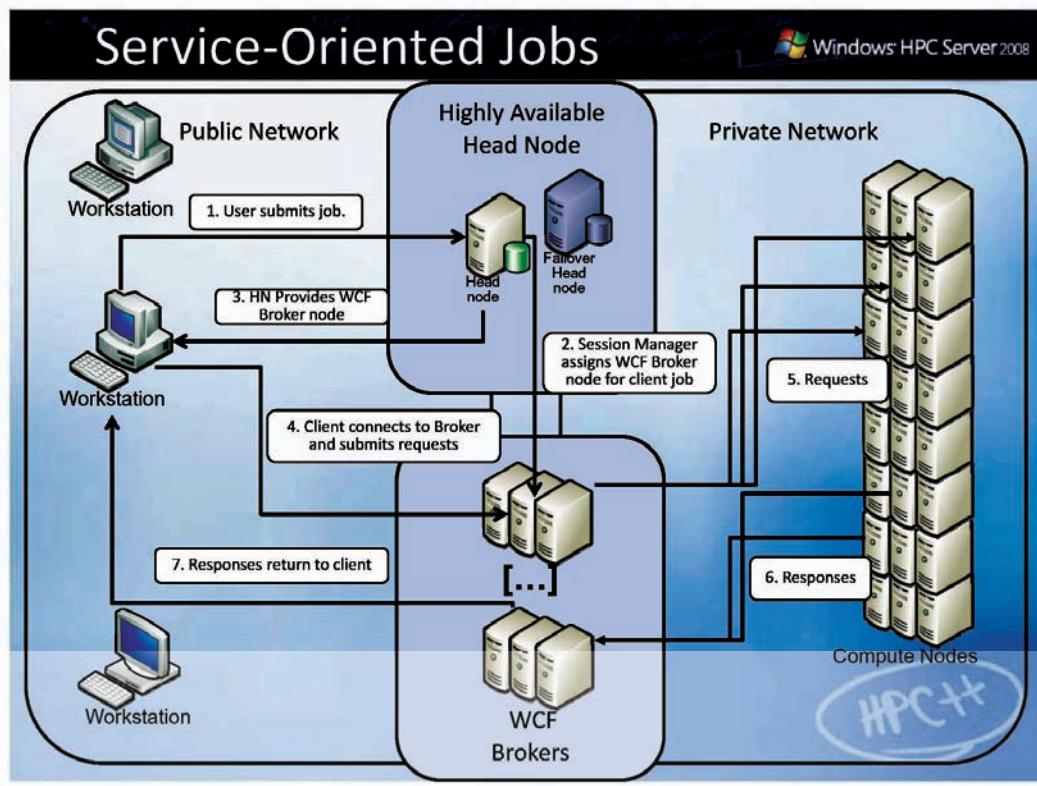
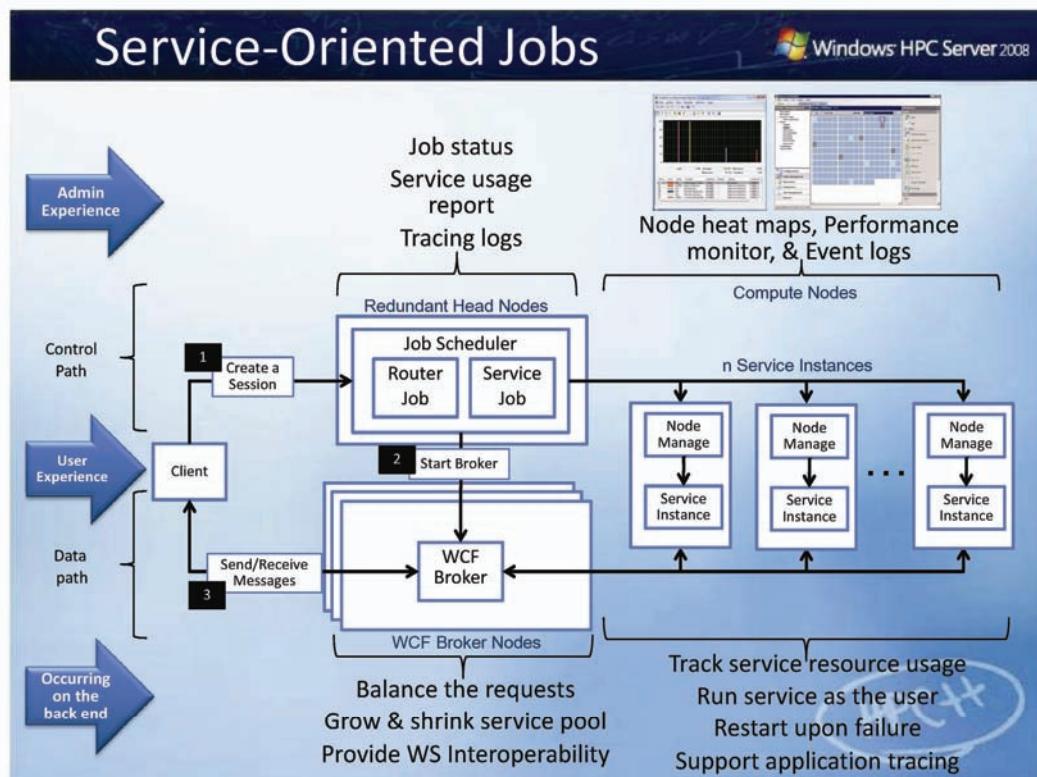
- Resource allocation
- Process Launching
- Resource usage tracking
- Integrated MPI execution
- Integrated Security

WCF Brokers

- WS Virtual Endpoint Reference
- Request load balancing
- Integrated Service activation
- Service life time management
- Integrated WCF Tracing

App.exe + Service (DLL)

HPC++



Job Templates

Windows HPC Server 2008

What are they?

- Encapsulates application resource requirements, job prioritization and resource allocation policies
- Created by the admin, ACL'ed to a set of users so that the resources can be delivered to the users in ways that meet the organization's business priorities

How to Define?

- Define the Value range: defines the range of values that the term can take. The value range can be empty. If the value provided at the job submission is not in the defined range, the job is rejected.
- Provide a default value: if a profile defines the default value, the user is not required to specify the job term

Job Templates

Windows HPC Server 2008

1 Admin creates resource partitions by creating job submission policies

Example Policies

Admission control	Descriptions	Definitions
Runtime to be mandatory	A supercomputing center wanting to enforce the runtime for all the jobs	Profile: default Runtime: required Default: none Users: All Profile LOB1: Users: user1, user2 Priority: normal, Select: "sas && ib && processorspeed > 2000000" Uniform: switchId Range: N/A Profile LOB2: Users: user3, user4 Askednodes:host2 host3 host 4 Profile PowerUser: Users: userA Askednodes: All Priority: Highest Range:
Multiple Line of Businesses (LOBs) sharing a cluster	Admin would like to apportion resources to different nodes	
Power user job priority	Power user userA can use all the nodes in the cluster	

2 Users submit using different templates

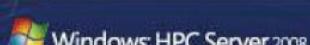
Five New Scheduling Policies



- Resource Matching
 - Job submit /nodegroup:appX myapp.exe
- Job Admission Policies via Templates
 - Job submit /template:groupY myapp.exe
- Multi-Level Processor Allocation
 - job submit /numsockets:4-8 myapp.exe
- Adaptive Allocation (Grow/Shrink)
 - job submit myapp.exe
- Preemption
 - Cluscfg PreemptionEnabled=true



Scenario: Placement via Job Context



node grouping, job templates, filters

Application Aware

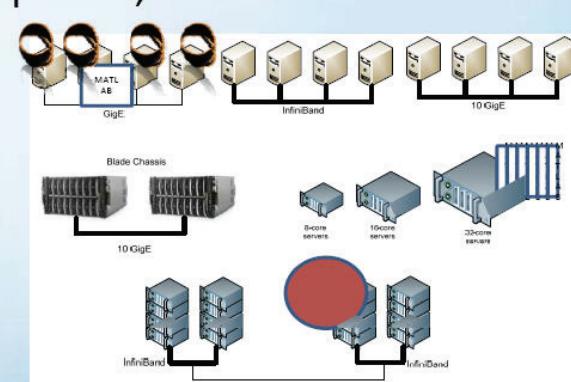
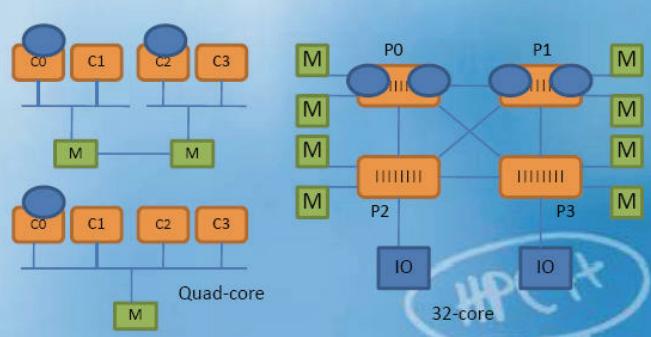
- [MATLAB icon] An ISV application (requires Nodes where the application is installed)

Capacity Aware

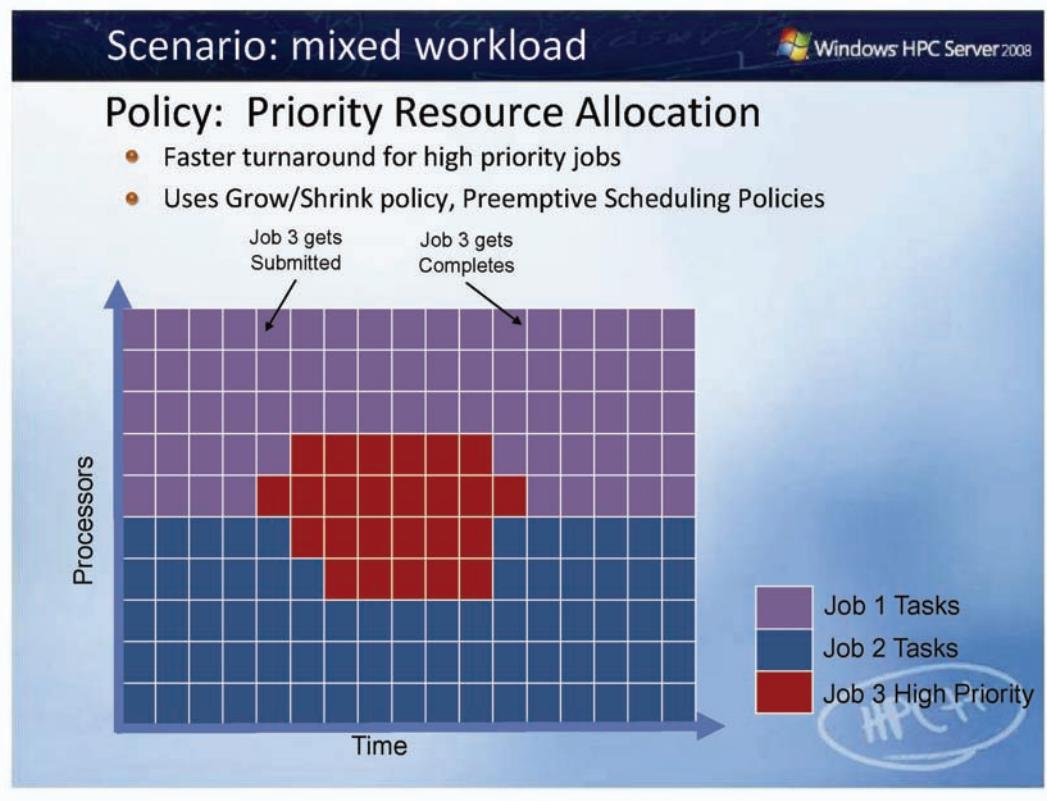
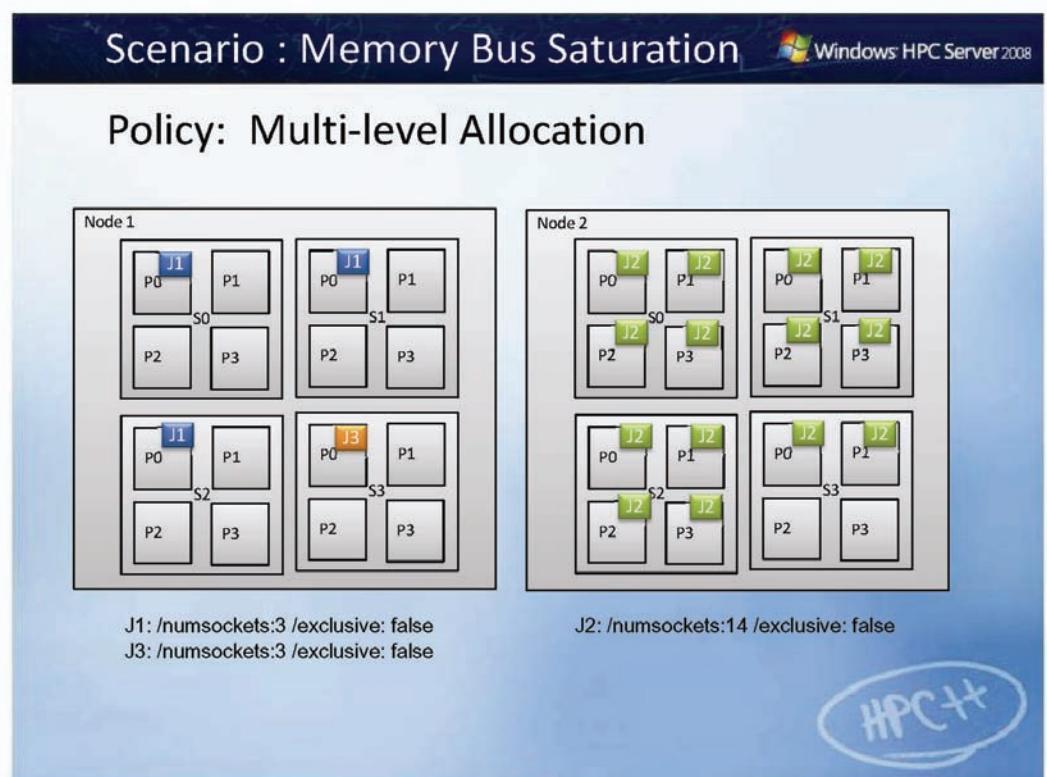
- [CPU icon] Multi-threaded application (requires machine with many Cores)
- [Red circle icon] A big model (requires Large memory machines)

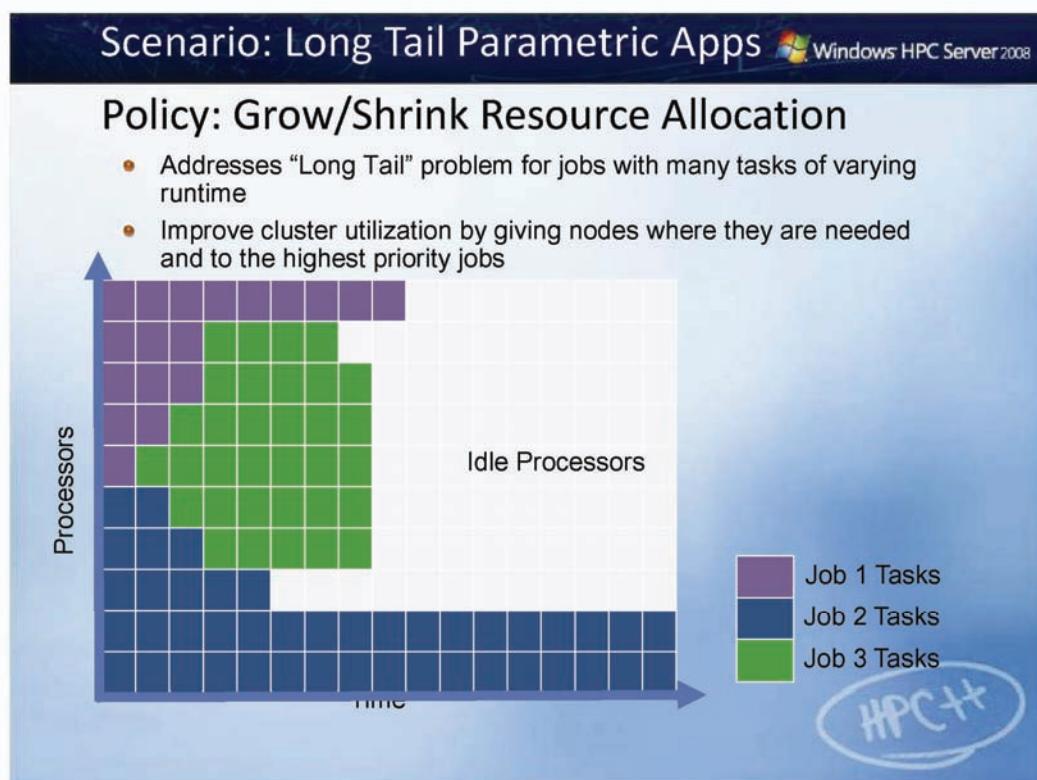
Numa Aware

- [Car icon] 4-way Structural Analysis MPI Job
- [Blue circles connected by arrows icon]

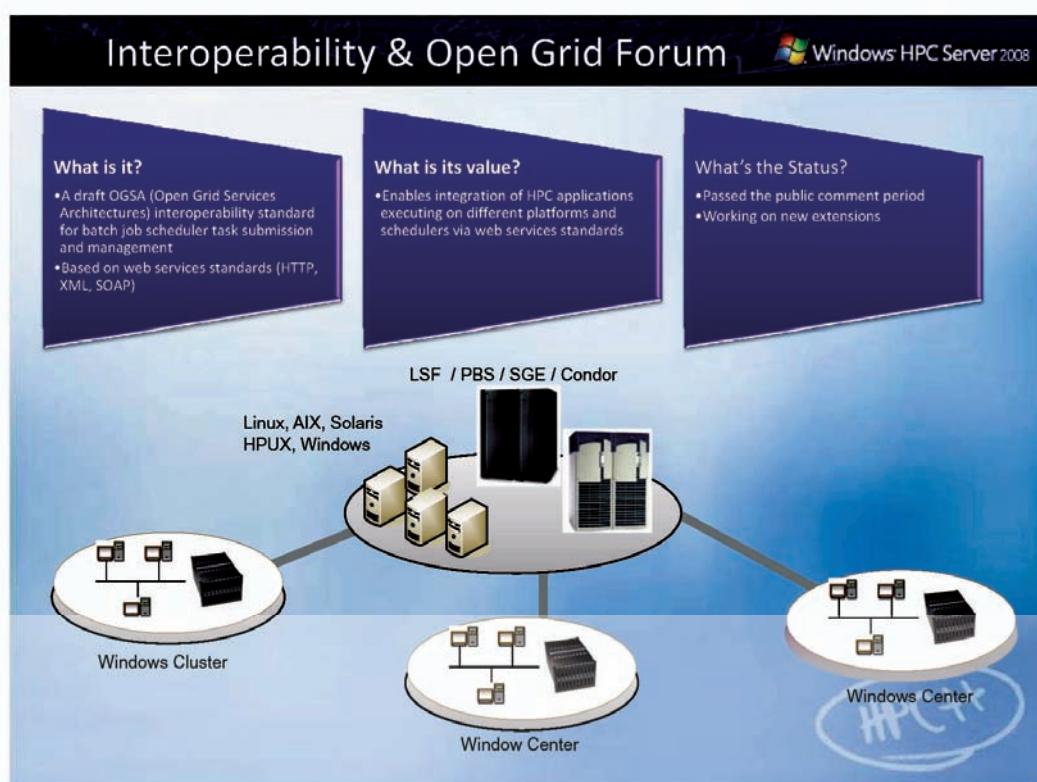






Job Policies Summary 

Features	Description	Version
Priority FCFS	Jobs are placed in the queue based on priority & submission Time. Jobs are in the queue on a First-Come First-Serve basis in the order they are submitted, except that ALL Highest Priority jobs are ahead of all Above Normal Priority jobs which are ahead of all Normal Priority jobs etc . . .	Window CCS 2003 Windows HPC 2008
Backfilling	Resources are reserved based on resource allocations. If the Job Scheduler identifies open windows for resources within the reserved timelines, it can backfill by selecting smaller jobs for execution within this window	Window CCS 2003 Windows HPC 2008
Exclusive Scheduling	When a job is “Exclusive” no other job can run on a node with that job. When a task is “Exclusive” no other task can run on a node with that task.	Window CCS 2003 Windows HPC 2008
Resource Match-making	Schedule against admin defined groups of nodes. We allow scheduling based on a limited amount of hardware properties and ordering by different hardware properties.	Windows HPC 2008
Multi-level Compute Resource Allocation	The Job Scheduler optimally places memory-intensive jobs to avoid contention of memory, delivering maximum and predictable application performance.	Windows HPC 2008
Preemption	Pre-emption occurs when high-priority jobs take resources away from lower-priority jobs which are already running.	Windows HPC 2008
Grow & Shrink Scheduling	A job may not have “uneven resource requirements” it may just need 100 resources to run 1000 tasks. With Grow Shrink, the job scheduler can give it 1 resource to get it started, and then additional resources as they become available. Once it has 100 resources and less than 100 tasks, job scheduler can begin taking away the un-needed resources.	Windows HPC 2008



Job Scheduling Summary

Windows HPC Server 2008

Tools	Windows CCS 2003	Windows HPC 2008
End User Tools	1. CLI 2. Job Manager UI	1. Powershell CLI 2. New Job Manager UI with built-in parametric support & custom job filtering
Admin Tools	1. Cluscfg for configuration 2. clusrun	1. Configuration UI 2. Enhanced clusrun
Developer Tools	1. C# API 2. COM API	1. WCF Integration 2. Scalable API with Rowset and eventing support 3. Standard Job Submission Interface (HPC Profile)
Scheduling Policies	1. FCFS 2. Backfill 3. Exclusive Scheduling 4. License scheduling	1. Resource matchmaking 2. Job Template 3. Multi-level processor allocation 4. Adaptive allocation / grow & shrink 5. Preemption



Thank you!



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Backup



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Panasas / pNFS / OSD Roadmap

